

where expectation is with respect to  $\tilde{\mathbf{w}}(t) = \mathbf{w}(t+1)$ , and note that

$$\begin{aligned} V_t(\mathbf{x}, \mathbf{y}) &= \max_{\mathbf{u} \in U_t(\mathbf{x})} E [g_t(\mathbf{x}, \mathbf{u}, \mathbf{y}) + V_{t+1}(\mathbf{f}_t(\mathbf{x}, \mathbf{u}, \mathbf{y}), \tilde{\mathbf{w}}(t))] \\ &= \max_{\mathbf{u} \in U_t(\mathbf{x})} \{g_t(\mathbf{x}, \mathbf{u}, \mathbf{y}) + E[V_{t+1}(\mathbf{f}_t(\mathbf{x}, \mathbf{u}, \mathbf{y}), \tilde{\mathbf{w}}(t))]\} \\ &= \max_{\mathbf{u} \in U_t(\mathbf{x})} \{g_t(\mathbf{x}, \mathbf{u}, \mathbf{y}) + G_{t+1}(\mathbf{f}_t(\mathbf{x}, \mathbf{u}, \mathbf{y}))\}. \end{aligned}$$

Now substituting  $\mathbf{y} = \mathbf{w}(t)$  and taking expectations with respect to  $\mathbf{w}(t)$  on both sides above we obtain

$$G_t(\mathbf{x}) = E \left[ \max_{\mathbf{u} \in U_t(\mathbf{x})} \{g_t(\mathbf{x}, \mathbf{u}, \mathbf{w}(t)) + G_{t+1}(\mathbf{f}_t(\mathbf{x}, \mathbf{u}, \mathbf{w}(t)))\} \right],$$

which gives us a recursion exactly of the form (D.3).

Note, however, that by using this transformation we have reduced the original dynamic programming recursion from one with a state space  $S_t \times W_t$  to one with only a state space of  $S_t$ . The function  $G_t(\mathbf{x})$  has a similar interpretation as  $V_t(\mathbf{x}, \mathbf{y})$  for this reduced state—namely, it is the optimal expected reward-to-go from time  $t$  onward given we are in the reduced state  $\mathbf{x}(t)$  at time  $t$ , where  $\mathbf{y} = \mathbf{w}(t)$  still uncertain (recall  $G_t(\mathbf{x}) = E[V_t(\mathbf{x}, \mathbf{w}(t))]$ ). Indeed, one can think of this new recursion as propagating the system in two stages: first, the state  $\mathbf{x}$  is realized but  $\mathbf{y}$  remains uncertain. We measure the optimal expected reward at this point, yielding  $G_t(\mathbf{x})$ . Then the value  $\mathbf{y} = \mathbf{w}(t)$  is realized, and we make our optimal decision. This takes us to a new state  $\mathbf{x}(t+1)$ , and the process repeats. Finally, note that this reduced-form recursion results in an optimization step of the form  $E[\max\{\ \}]$  rather than the  $\max E[\{\ \}]$  found in traditional dynamic programming formulations.

Here's a typical example of how this transformation arises in RM. Suppose  $x(t)$  is a scalar capacity,  $y(t)$  is the revenue of the request in period  $t$ , and  $u(t) = 1$  if we decide to accept a request and zero otherwise. So the reward function is simply

$$y(t)u(t).$$

Capacity evolves according to the system equation

$$x(t+1) = x(t) - u(t),$$

and the revenue is driven by a random process

$$y(t+1) = w(t+1).$$

Formulated in traditional terms, we obtain

$$V_t(x, y) = \max_{u \in \{0,1\}} E [yu + V_{t+1}(x - u, w(t))].$$

However, with the transformation above, we can rewrite this in observable-disturbance form as

$$G_t(x) = E \left[ \max_{u \in \{0,1\}} \{w(t)u + G_{t+1}(x - u)\} \right].$$

Since most dynamic programs in RM are of this observable-disturbance form, we typically use the simpler  $E[\max\{\ \}]$  rather than the traditional  $\max E[\{\ \}]$  form.